

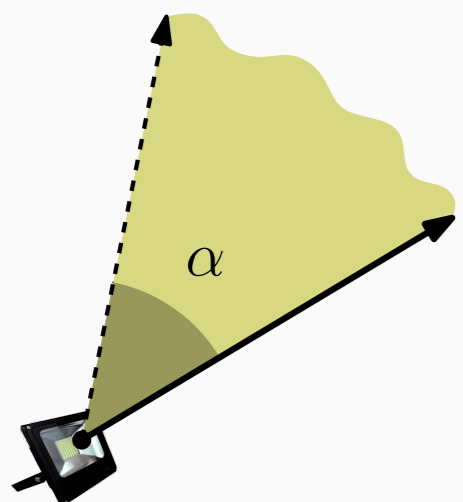
THE VORONOI DIAGRAM OF ROTATING RAYS

WITH APPLICATIONS TO FLOODLIGHT ILLUMINATION

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Brocard Illumination Problem

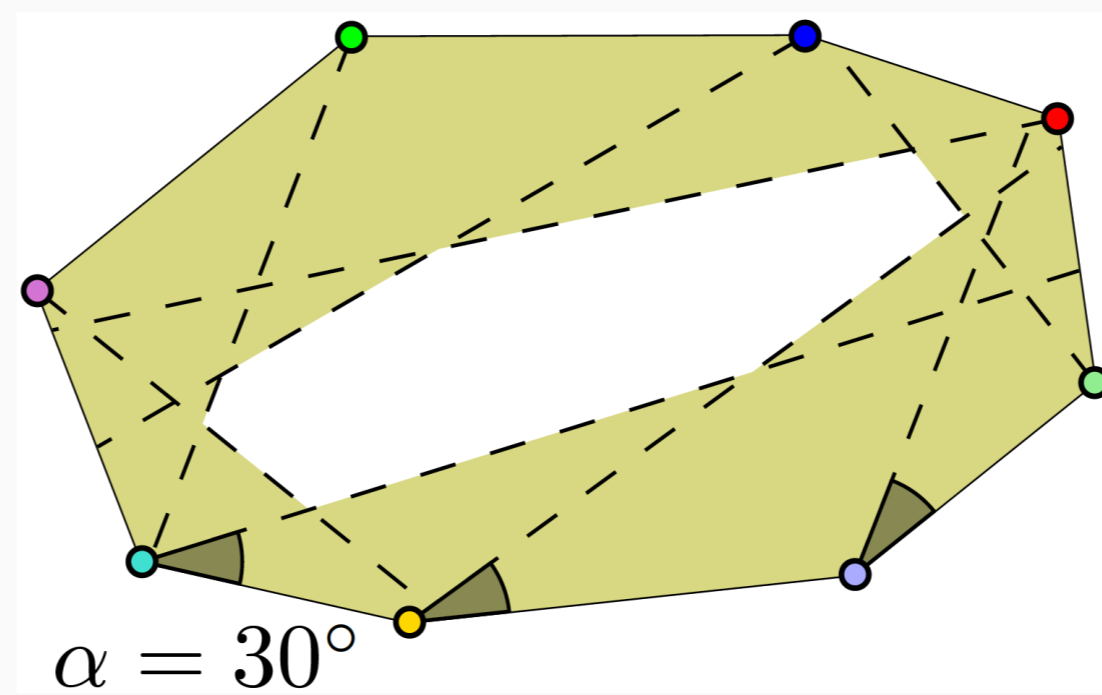
α -floodlight



Problem

Given a polygon P with an α -floodlight aligned at each edge, what is the **minimum angle** α^* needed to illuminate P ?

α^* is the **Brocard angle** of P .

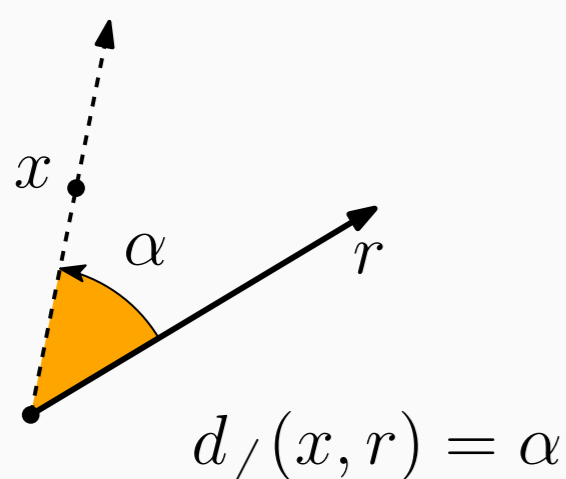


Approach

We define the **Rotating Rays Voronoi Diagram (RVD)**. The sites are **rays** and the distance is the **angular distance**. The Brocard angle is realized at a vertex of the RVD.

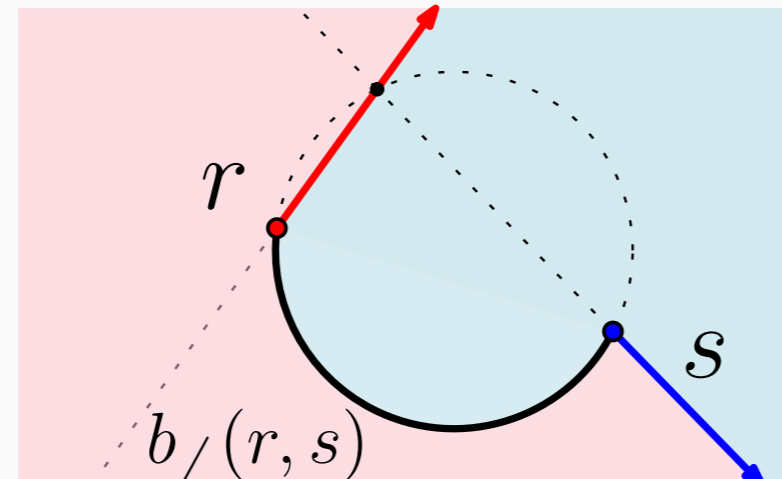
Angular Distance

The **angular distance** of a ray r to point x : $d_{\angle}(x, r)$.

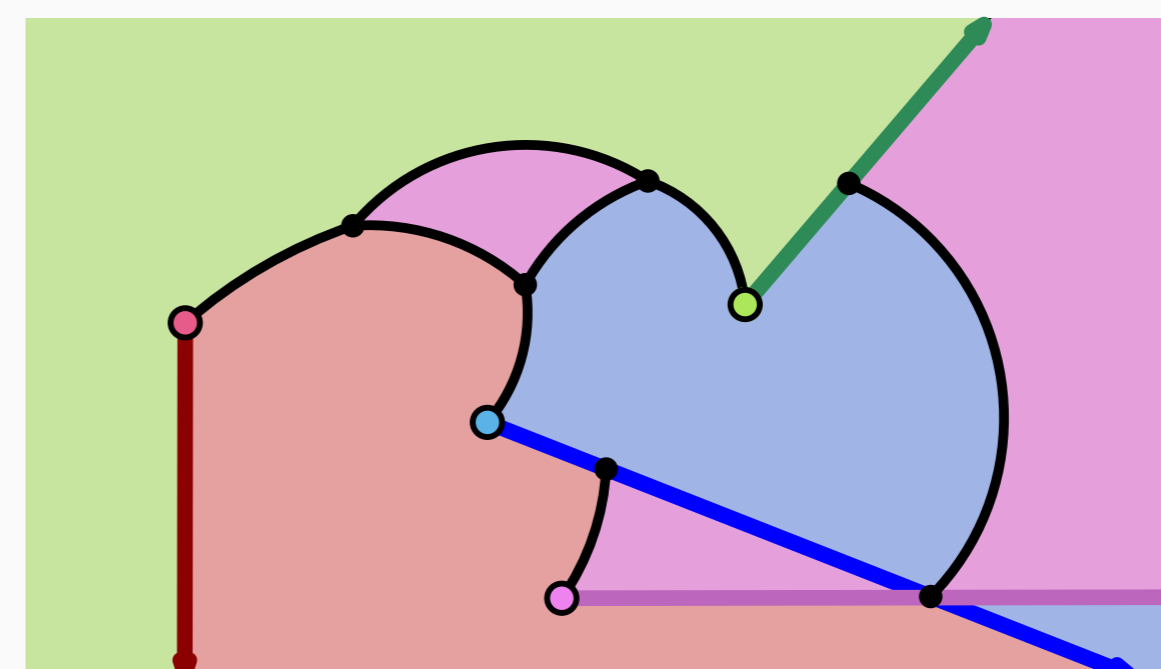


Angular Bisector

The **angular bisector** of two rays r and s : $b_{\angle}(r, s)$.

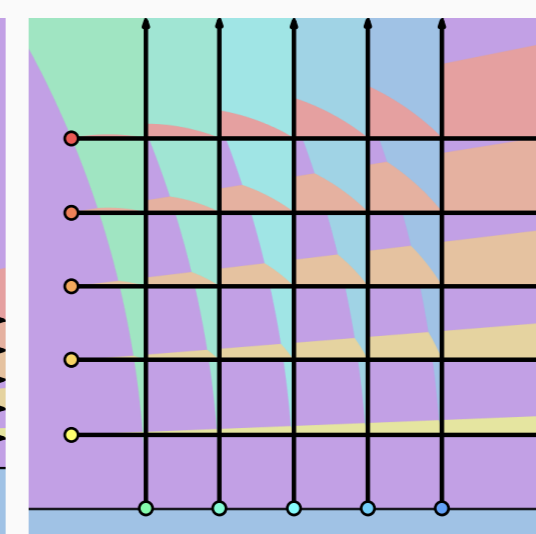
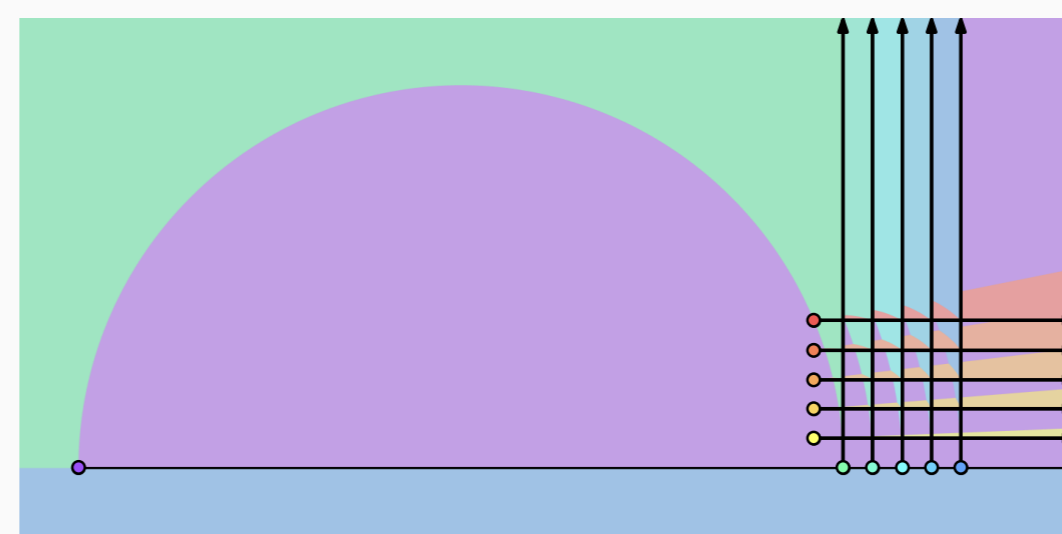
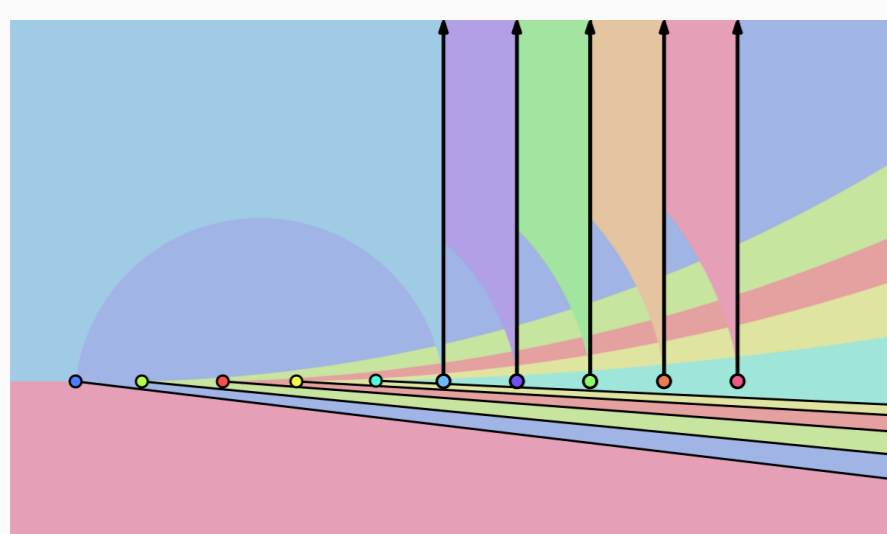


Rotating Rays Voronoi Diagram



The **Voronoi region** of $r \in \mathcal{S}$ is the locus of points closer to r . The **Rotating Rays Voronoi Diagram** of \mathcal{S} is the subdivision of \mathbb{R}^2 in Voronoi regions.

Rotating Rays Voronoi Diagram in \mathbb{R}^2



Theorem

$\text{RVD}(\mathcal{S})$ has $\Omega(n^2)$ **worst case complexity** even if the rays are non-intersecting.

Theorem

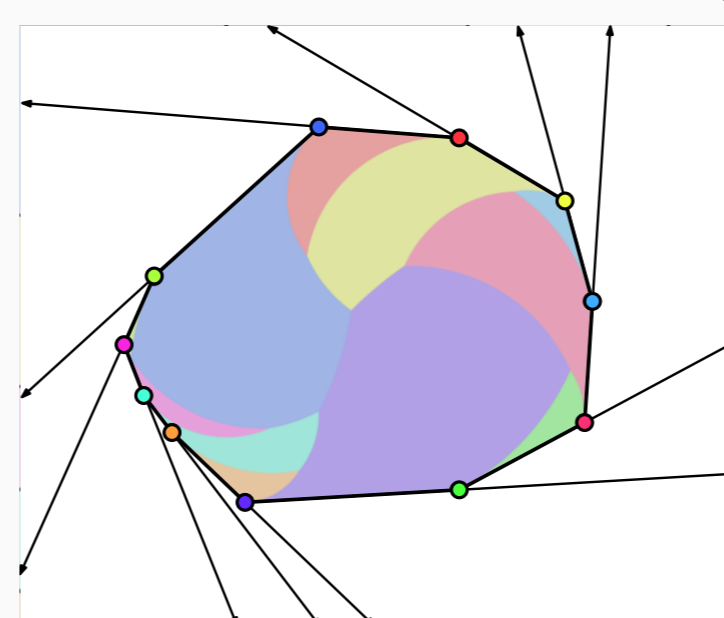
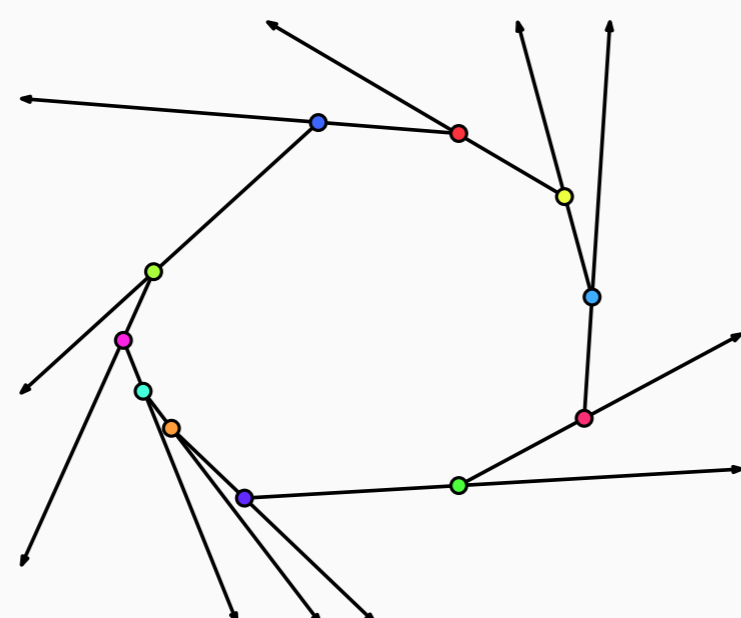
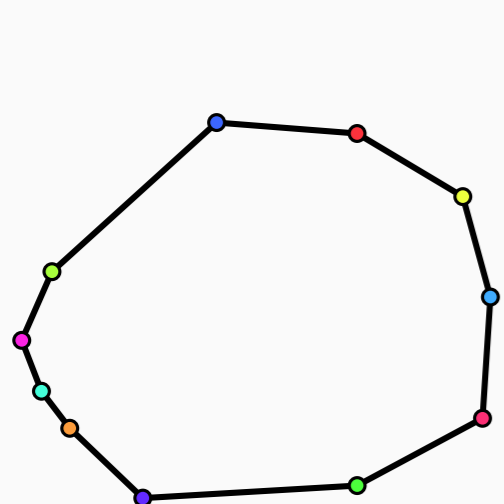
A single Voronoi region of $\text{RVD}(\mathcal{S})$ can have $\Theta(n^2)$ **complexity** in the worst case.

Theorem

$\text{RVD}(\mathcal{S})$ has $O(n^{2+\epsilon})$ **complexity** and it can be constructed in $O(n^{2+\epsilon})$ **time**.

Rotating Rays Voronoi Diagram of a Convex Polygon

Input: Convex polygon P **Obtain:** Set of rays \mathcal{S}_P **Output:** Diagram $\text{PRVD}(\mathcal{S}_P)$



Algorithm outline

1. Partition \mathcal{S}_P in 4 sets.
2. Construct each of the 4 diagrams.
3. Merge the 4 diagrams.

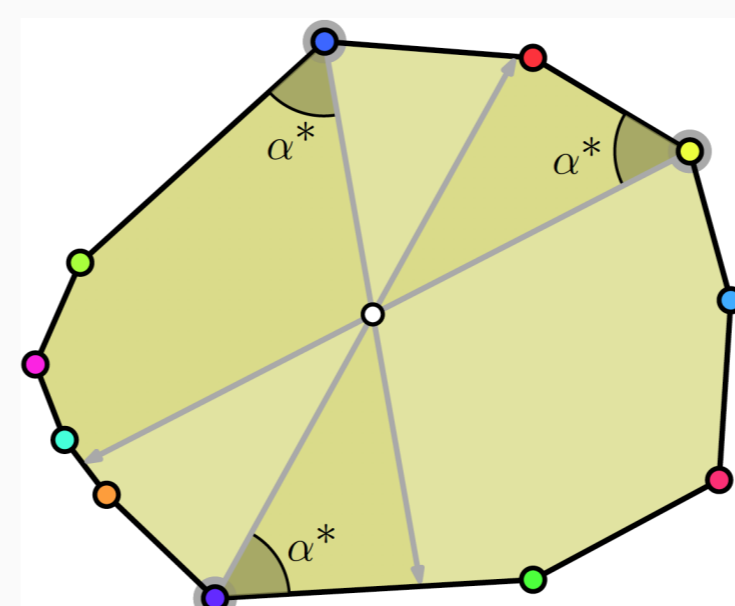
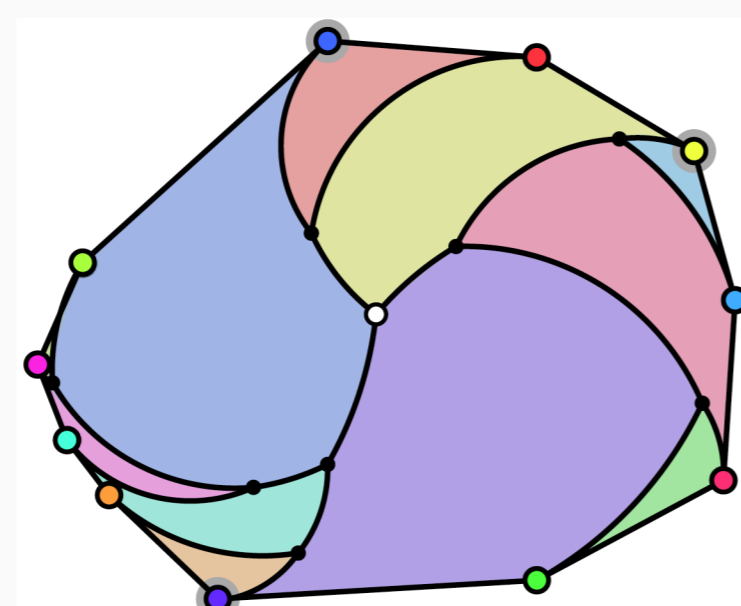
Theorem

$\text{PRVD}(\mathcal{S}_P)$ can be constructed in $\Theta(n)$ **time**.

Finding the Brocard angle

Brocard angle of polygon P

1. Create $\text{PRVD}(\mathcal{S}_P)$
 2. Traverse $\text{PRVD}(\mathcal{S}_P)$
- **Three α^* -floodlights suffice**
 - $\alpha^* \in (0, \pi/2 - \pi/n]$



Brocard angle of \mathcal{S} in \mathbb{R}^2

1. Create $\text{RVD}(\mathcal{S})$
 2. Traverse $\text{RVD}(\mathcal{S})$
- α^* realized at vertex of $\text{RVD}(\mathcal{S})$
 - $\alpha^* \in (2\pi/n, 2\pi)$